

Optimal Finite Blocklength Code Design on the Binary Erasure Channel

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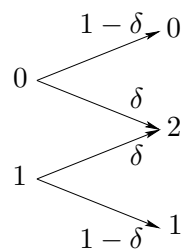
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1 Introduction

In traditional coding theory, it is the goal to find good codes that operate close to the ultimate limit of the **channel capacity** as introduced by Shannon. In this work we would like to break away from these traditional simplifications and instead focus on an optimal (i.e., **minimum error probability**) design of codes for a certain given channel and a given fixed blocklength. Since for very short blocklength, it is not realistic to transmit large quantities of information, we start by looking at codes with only a few codewords, so called **ultra-small block-codes**. We introduce a new class of codes, called **fair weak flip codes** (the number of zeros and ones are almost equal under certain conditions on the blocklength n) and prove that they have beautiful **quasi-linear** properties.

2 Channel Models

We consider the **binary erasure channel (BEC)** with erasure probability δ :



3 Optimal Code Design

Define the following fair weak flip codes:

$$\mathcal{C}_{\text{BEC}}^{(5,10)*} \triangleq \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 \end{pmatrix}$$

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$$\mathcal{C}_{\text{BEC}}^{(6,10)*} \triangleq \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 0 \end{pmatrix}$$

Theorem 1. For a BEC and for any n being a multiple of 10, an **optimal codebook** with $M = 5$ or $M = 6$ codewords is

$$\left(\mathcal{C}_{\text{BEC}}^{(M,10)*} \mathcal{C}_{\text{BEC}}^{(M,10)*} \dots \mathcal{C}_{\text{BEC}}^{(M,10)*} \right)_{M \times (n \bmod 10 = 0)}$$

4 Optimal Exact Error Probability

Optimal Performance of BEC with $M = 4$ and $\delta = 0.3$

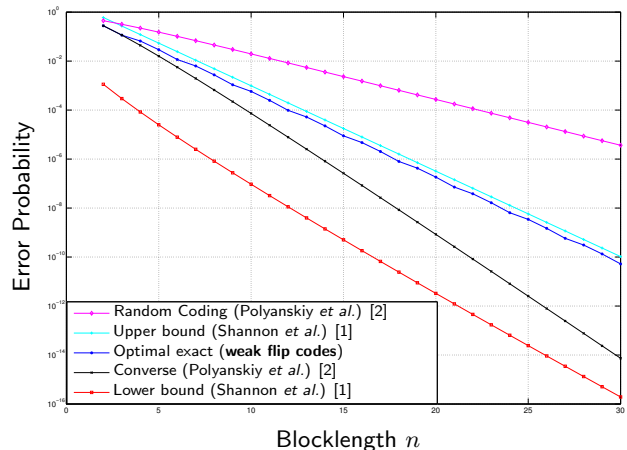


Figure 1: A comparison of some known bounds and the exact error probability of the globally optimal code with $M = 4$ codewords on a BEC.

References

- [1] C. E. Shannon, R. G. Gallager, and E. R. Berlekamp, "Lower bounds to error probability for coding on discrete memoryless channels," *Inform. Contr.*, pp. 522–552, May 1967, part II.
- [2] Y. Polyanskiy, H. V. Poor, and S. Verdú, "Channel coding rate in the finite blocklength regime," *IEEE Trans. Inf. Theory*, vol. 56, no. 5, pp. 2307–2359, May 2010.